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# Exploring Go's Prowess in Immersive Database Development: immudb as a Case Study



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- 1. immudb: The Immutable Database
- 2. Why Golang?
- 3. The Garbage Collector
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- **immudb** is a database with built-in cryptographic proof and verification
- It operates as a key-value store, SQL or document oriented database.
- Engineered for scalability, immudb can handle millions of transactions per second.

## **Project History**

- Started in **2020**
- **4785** commits, **~70** contributors
- 47 releases (latest version is v1.9.5)
- ~146k lines of code (Golang only)



## Applications of an immutable database



## Immutable append only logs

• An append-only log file is a data structure where new entries are added sequentially at the end, and existing entries cannot be modified or deleted.

Append Only

### immudb: the Merkle Tree



## The inclusion proof

• It proofs a given leaf's membership in the tree



## The consistency proof

 It proofs that a merkle tree (root hash) is a superset of a previous merkle tree



How immutability is proved to clients?

- 1. **Root Hash Tracking**: clients retain the most recent known Merkle tree root hash *H*;
- 2. **Proofs for Integrity**: when fetching or setting a new key-value entry *E*, clients can optionally request:
  - a. a *consistency proof* between *H* and the latest database state *H*';
  - b. an *inclusion proof* confirming that *E* belongs to the Merkle tree represented by the current root hash *H*'.

## Why Golang?

- Extensive standard library, including cryptographic support
- Native concurrency features (goroutines, channels and more)
- Successful examples of adoption in large database and storage projects (CockroachDB, BadgerDB)
- Tools for reducing GC load: sync.Pool, unsafe package, etc.

### **Rich Standard Library**

```
func (t *HTree) BuildWith(digests [][sha256.Size]byte) error {
       . . .
       b := [1 + 2*sha256.Size]byte{NodePrefix}
       for w > 1 {
               wn := 0
               for i := 0; i+1 < w; i += 2 {
                       copy(b[1:], t.levels[1][i][:])
                       copy(b[1+sha256.Size:], t.levels[1][i+1][:])
                       t.levels[1+1][wn] = sha256.Sum256(b[:])
                       wn++
               if w%2 == 1 {
                       t.levels[1+1][wn] = t.levels[1][w-1]
                       wn++
               }
               1++
               w = wn
ł
```

## Goroutines

```
resCh := make(chan appendableResult)
go func() {
      offsets, err := s.appendValuesToAnyVLog()
      if err != nil {
            resCh <- appendableResult{nil, err}</pre>
      }
      resCh <- appendableResult{offsets, nil}</pre>
}()
err = tx.BuildHashTree()
. . .
valuesRes := <-resCh</pre>
. . .
```

## Channels

- Channels enable safe communication between goroutines, avoiding explicit locks and simplifying concurrent code.
- For simple local data sharing, channels might add unnecessary complexity, as opposed to direct data passing or using mutexes.
- Circular waiting patterns, where goroutines are blocked waiting on each other, can lead to deadlocks.



## Channels in the WatchersHub component



```
WatchersHub usage
```



## WatchersHub: implementation

```
type WatchersHub struct {
    wPoints map[uint64]waitPoint
    value uint64
    ...
    mtx sync.Mutex
}
```

type waitPoint struct {
 v uint64
 ch chan struct{}
 count int
}

## WatchersHub: the WaitFor method

```
func (w *WatchersHub) WaitFor(ctx context.Context, v uint64) error {
       if w.value >= v {
               return nil
       wp, waiting := w.wPoints[v]
       if !waiting {
               wp = waitingPoint{v: v, ch: make(chan struct{})}
        . . .
       cancelled := false
       select {
               case <-wp.ch:</pre>
                       break
               case <-ctx.Done():</pre>
                       cancelled = true
```



## WatchersHub: the Advance method

```
func (w *WatchersHub) Advance(v uint64) error {
      . . .
      for i := s.value+1; i <= v; i++ {</pre>
            wp, waiting := w.wPoints[i]
            if waiting {
                  close(wp.ch)
                   . . .
                  delete(w.wPoints, i)
             }
      w.value = v
      . . .
}
```

- Use sync.Pool to reduce the number of total allocations
- Prefer value types over pointers when feasible
- Eliminate unnecessary copying

The mark and sweep algorithm

- It walks the program's object graph to identify objects that are in use by the program
- It uses a two phase technique, called **mark**-and-**sweep**.



# Mark and sweep (MARK)





# Mark and sweep (SWEEP)



The GC Cycle

- GC cycles pause time increases as we allocate more and more objects
- It is not affected by the total amount of bytes allocated

## Reducing the number of allocations with sync.Pool

```
type Row struct {
    . . .
var rowsPool = sync.Pool {
     New: func() any {
           return &sql.Row{...}
s := rowsPool.Get().(*sql.Row)
s.reset()
... // do something with r
rowsPool.Put(s) // release r back to the pool
```

## Prefer values to pointers



## Cache





type	cacheShard struct {
	offsets map[string]uint32
	buf []byte
}	
func	(c *cacheShard) <b>Put(key, value []byte) error {</b>
}	
func	(c *cacheShard) Get(key []byte) ([]byte, error) {
}	



## Conclusions

- Golang simplifies complex tasks, from cryptography to concurrency, with minimal dependencies.
- While GC can introduce delays, there are several options to minimize its impact.
- Projects like immudb showcase its effectiveness in real-world, performance critical applications.

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https://codenotary.com

https://immudb.io

https://github.com/codenotary/immudb



# Thank You For Your Time!



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